

Formalising the classification of groups of order pq

Peiran Wu

University of St Andrews

25 September 2024

About Lean

- ▶ Open-source functional programming language
- ▶ Interactive theorem prover – used to confirm the correctness of mathematical proofs
- ▶ Accompanied by "Mathlib", a community-maintained open-source mathematical library
- ▶ Google DeepMind's AlphaProof (together with AlphaGeometry 2)
- ▶ Lots of gaps in the group theory section of Mathlib
- ▶ Demo – how to read Lean code

Groups of order pq

- ▶ Last year, there was a study group on Lean in St Andrews.
- ▶ As an exercise, I wanted to classify the groups of order 4, but someone beat me to it.
- ▶ Just before Christmas, I formalised the classification of groups of order 6 (140 LoC, not golfed) and felt it wouldn't be too hard to generalise.
- ▶ Scott Harper and I started working together on the classification of groups of order pq .

Let p and q be positive prime numbers with $p \leq q$. Let G be a group of order pq . Then exactly one of the following holds:

- (1) $G \cong C_{pq}$.
- (2) $p = q$ and $G \cong C_p \times C_p$.
- (3) $p \mid q - 1$, G is non-abelian, and $G \cong C_q \rtimes C_p$.

The statement

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- (3) $p \mid q - 1$, G is non-abelian, and $G \cong C_q \rtimes C_p$.

A few things we had to take into consideration:

- (a) Implicit: the semidirect product involves a non-trivial homomorphism $C_p \rightarrow \text{Aut } C_q$.
- (b) Implicit: the choice of this homomorphism does not matter.
- (c) Implicit: such a homomorphism exists.
- (d) The natural-language statement has the usual form of a classification result. This may not be user-friendly in Lean.

Rephrasing

We ended up choosing the following main theorems plus some corollaries:

Let p and q be positive prime numbers with $p < q$.

- (i) If G is a non-cyclic group of order p^2 , then G is isomorphic to $C_p \times C_p$.*
- (ii) If G is a group of order pq and $p \nmid q - 1$, then G is cyclic.*
- (iii) If $p \mid q - 1$, then there exists a non-cyclic group of order pq .*
- (iv) If G is a non-cyclic group of order pq and $\varphi : C_p \rightarrow \text{Aut } C_q$ is any non-trivial homomorphism of groups, then G is isomorphic to $C_q \rtimes_{\varphi} C_p$.*

Finiteness and cardinality

“Let G be a finite group of cardinality n .”

```
1 (G : Type*) [Group G] [Finite G] (n : N)
```

```
(h : Nat.card G = n)
```

```
2 (G : Type*) [Group G] [Fintype G] (n : N)
```

```
(h : Fintype.card G = n)
```

```
class inductive Finite (α : Sort*) : Prop
```

```
| intro {n : N} : α ≈ Fin n → Finite _
```

```
class Fintype (α : Type*) where
```

```
elems : Finset α
```

```
complete : ∀ x : α, x ∈ elems
```

```
def Nat.card (α : Type*) : N := ...
```

```
def Fintype.card (α : Type*) [Fintype α] : N := ...
```

Finite vs Fintype

```
instance Finite.of_fintype (a : Type*) [Fintype a] : Finite a
```

```
noncomputable def Fintype.ofFinite (a : Type*) [Finite a] : Fintype a
```

Fintype is meant to be used for computable definitions and in general requires more work (without resorting to `classical`). For example, we initially had to prove the following.

```
instance MulEquiv.fintype (α β : Type*) [DecidableEq α] [DecidableEq β]  
  [Mul α] [Mul β] [Fintype α] [Fintype β] : Fintype (α ≈* β) where  
  ...
```

```
instance Fintype.decidableEqMulEquivFintype (α β : Type*)  
  [DecidableEq β] [Fintype α] [Mul α] [Mul β] : DecidableEq (α ≈* β) :=  
fun a b => decidable_of_iff ((a : α → β) = b)  
  (Injective.eq_iff DFunLike.coe_injective)
```

Switching from `Fintype` to `Finite`

Previously, Lagrange's Theorem:

```
theorem Subgroup.card_subgroup_dvd_card
  {α : Type u_1} [Group α] (s : Subgroup α) [Fintype α] [Fintype s] :
  Fintype.card ↑s | Fintype.card α :=
  ...
```

The `[Fintype s]` is needed unless we use `classical`.

As of June:

```
theorem Subgroup.card_subgroup_dvd_card
  {α : Type u_1} [Group α] (s : Subgroup α) :
  Nat.card ↑s | Nat.card α :=
  ...
```

This also covers infinite groups.

Homomorphisms

`variable (G1 G2 : Type*) [Group G1] [Group G2]`

The type of group homomorphisms from G_1 to G_2 is `MonoidHom G1 G2`, notation $G_1 \rightarrow^* G_2$.

Similarly, `MulEquiv G1 G2`, notation $G_1 \simeq^* G_2$, is the type of group isomorphisms (or in fact, semigroup isomorphisms).

An object of type $G_1 \simeq^* G_2$ is a concrete isomorphism. When we only care that the groups are isomorphic, we can write `Nonempty (G1 \simeq^* G2)`.

How to say a homomorphism $\varphi : G_1 \rightarrow^* G_2$ is non-trivial? At first we came up with `φ .ker \neq T` and `φ .range \neq \perp` . It turned out we can just write `$\varphi \neq 1$` .

Cyclic groups

In Mathlib, `ZMod n` is the type of integers modulo n for non-zero n and we know that it is a commutative ring and it is cyclic as a group.

```
instance ZMod.commRing (n : ℕ) : CommRing (ZMod n)
```

```
instance ZMod.instIsAddCyclic (n : ℕ) : IsAddCyclic (ZMod n)
```

Unfortunately for us, additive notation is used for its group structure. We replace this with multiplicative notation in a new type:

```
abbrev MulZMod (n : ℕ) := Multiplicative (ZMod n)
```

```
instance isCyclic_multiplicative {α : Type u} [AddGroup α] [IsAddCyclic α] :  
  IsCyclic (Multiplicative α) :=
```

```
...
```

Stating (iv) in Lean

Let p and q be positive prime numbers with $p < q$. If G is a non-cyclic group of order pq and $\varphi : C_p \rightarrow \text{Aut } C_q$ is any non-trivial homomorphism of groups, then G is isomorphic to $C_q \rtimes_{\varphi} C_p$.

```
variable {p q : ℕ} (hp : p.Prime) (hq : q.Prime) (hpq : p < q)
```

```
variable {G : Type*} [Group G] [Finite G]
```

```
theorem mulEquiv_semidirectProduct_of_not_isCyclic_of_card
```

```
  (h : Nat.card G = p * q) (h' : ¬IsCyclic G)
```

```
  (φ : MulZMod p →* MulAut (MulZMod q)) (hφ : φ ≠ 1) :
```

```
  Nonempty (G ≃* MulZMod q ⋊[φ] MulZMod p) :=
```

```
  ...
```

Proof of (iv)

Let p and q be positive prime numbers with $p < q$. If G is a non-cyclic group of order pq and $\varphi : C_p \rightarrow \text{Aut } C_q$ is any non-trivial homomorphism of groups, then G is isomorphic to $C_q \rtimes_{\varphi} C_p$.

- ▶ Let G be a non-cyclic group of order pq . Let $\varphi : C_p \rightarrow \text{Aut } C_q$ be non-trivial.
- ▶ By Sylow's Theorems, there is a subgroup H of order q and a subgroup K of order p in G ; furthermore the number n of subgroups of order q in G is $1 \pmod q$.
- ▶ The index $|G : N_G(H)|$ is also n and divides p by a corollary of Lagrange's Theorem. Therefore H is normal in G .
- ▶ The set $HK = \{hk : h \in H, k \in K\}$ has cardinality $|H||K|/|H \cap K|$. Since H and K intersect trivially, HK has cardinality pq . Therefore HK is the biggest set (and hence subgroup) in G .
- ▶ K acts on H by conjugation, giving rise to some $\psi : K \rightarrow \text{Aut } H$ and some $G \simeq H \rtimes_{\psi} K$.
- ▶ Since there are instances of $H \simeq C_q$ and $K \simeq C_p$, and ψ is compatible with φ , we obtain some $G \simeq C_q \rtimes_{\varphi} C_p$ by a congruence lemma for semidirect products.

A key lemma

$\text{Aut } C_q$ is isomorphic to C_{q-1} .

(Hence if $p \mid q-1$, then there exists a unique non-trivial homomorphic image of C_p in $\text{Aut } C_q$; and if $p \nmid q-1$, then any $C_p \rightarrow^* \text{Aut } C_q$ is trivial.)

We laid out many options for a proof assuming Mathlib and chose the one with the least amount of mathematical content:

```
variable (p : N) [Fact (p.prime)]
```

```
def addEquivAddAutZMod : AddAut (ZMod p)  $\simeq^*$  (ZMod p)x where ...
```

```
def mulEquivMulAutMulZMod : MulAut (MulZMod p)  $\simeq^*$  (ZMod p)x :=  
  AddEquiv.toMultiplicative.mulEquiv.symm.trans <| addEquivAddAutZMod p
```

```
lemma mulAut_MulZMod_isCyclic : IsCyclic (MulAut (MulZMod p)) := ...
```

```
lemma card_mulAut_mulZMod :  
  Nat.card (MulAut (MulZMod p)) = p - 1 := ...
```

Concluding the project

- ▶ The code is completely `sorry`-free at about 1000 lines.
- ▶ We have learned a lot about Lean and Mathlib.
- ▶ We are in the process of refactoring, generalising our lemmas, and submitting pull requests to Mathlib in small chunks.
- ▶ We feel there is room for more automation, and more documentation and organisation.
- ▶ Many more theorems in group theory and lemmas about finite objects in general await formalisation. And maybe a computational library?