Card Shuffle Groups

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This is joint work with Binzhou Xia and Zhishuo Zhang.

7 Nov 2023, Seminar (Online)

Perfectly shuffle 2n cards

• Cut the deck in half:



• Perfectly interleave them:



Perform out-shuffles on a deck of 52 cards repeatedly.

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Question: What is the minimum number of times needed to return to the original order?

Answer: 8 times.

• Position x: 0 1 2 3 4 ··· 25 26 ··· 50 51



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Position x: 0
1
2
3
4
25
26
50
51
after O: 0
26
1
27
2
...
38
13
...
25
51



• Position x: 0 1 2 3 4 ... 25 26 ... 50 51 after O: 0 26 1 27 2 ... 38 13 ... 25 51 $x^{O}: 0$ 2 4 6 8 ... 50 1 ... 49 51



• Position x: 0 1 2 3 4 \cdots 25 26 \cdots 50 51 after O: 0 26 1 27 2 \cdots 38 13 \cdots 25 51 $x^{O}:$ 0 2 4 6 8 \cdots 50 1 \cdots 49 51

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$$(i + 26j)^O = 2i + j$$
 for $i \in \{0, ..., 25\}$
and $j \in \{0, 1\}$;



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$$x^{O} = (2x \mod{51}) \text{ for } x \in \{1, \dots, 50\};$$



Out-shuffle O

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 the order of O is the smallest positive integer t such that 2^t ≡ 1 (mod 51).



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• $(i + jn)^O = 2i + j$ for $i \in \{0, ..., n - 1\}$ and $j \in \{0, 1\}$;

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- $(i + jn)^O = 2i + j$ for $i \in \{0, ..., n 1\}$ and $j \in \{0, 1\}$;
- $x^{O} = (2\ell \mod 2n 1)$ for $x \in \{1, \dots, 2n 2\}$.

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To answer these questions, we first determine the parity of O and I.

Observation: for n = 4,

• Position x: 0 1 2 3 4 5 6 7
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$$x^{O}$$
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For general n, the order of the 2n cards after O is

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For general n, the order of the 2n cards after O is

$$(0, 2, 4, 6, \ldots, 2n - 2, 1, 3, 5, \ldots, 2n - 1),$$

and thus its inversion number is $1 + \cdots + n - 1 = n(n-1)/2$.

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- *I* is obtained by permutating the two piles and then performing *O*.
 (x^I = x^{(0,n)(1,n+1)…(n-1,2n-1)O} for all x ∈ {0,1,...,2n-1})

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- If *n* and *O* have the same parity, then *I* is even; otherwise *I* is odd \implies If $n \equiv 0$ or 3 (mod 4), then *I* is even; otherwise *I* is odd.
- Thus $\langle O, I \rangle \leq \operatorname{Alt}(2n) \iff n \equiv 0 \pmod{4}$.

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- out-shuffle and in-shuffle preserve the partition $\{0,7\}, \{1,6\}, \{2,5\}, \{3,4\}.$

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For a general *n*, out-shuffle and in-shuffle preserve the partition $\{0, 2n - 1\}, \{1, 2n - 2\}, \dots, \{n - 1, n\}.$

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- Alt(2*n*) is primitive \Longrightarrow Alt(2*n*) $\notin \langle O, I \rangle$

- A permutation group on a set Ω is said to be imprimitive if this group preserves a nontrivial partition of Ω; otherwise, it is said to be primitive.
- *O* and *I* preserve $\{0, 2n 1\}, \{1, 2n 2\}, \dots, \{n 1, n\}$ $\implies \langle O, I \rangle$ is imprimitive.
- Alt(2n) is primitive \Longrightarrow Alt(2n) $\leq \langle O, I \rangle \Longrightarrow O$ and I can't generate Alt(2n) or Sym(2n).

Diaconis-Graham-Kantor

Question: How many/Which orderings can be obtained by performing a sequence of the two shuffles?

[1] P. Diaconis, R. L. Graham and W. M. Kantor, The mathematics of perfect shuffles., *Adv. Appl. Math.*, 4 (1983), 175–196.

Wenying Zhu (BNU)

Card Shuffle Groups

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Answered by Diaconis, Graham and Kantor in 1983^[1].



Persi Diaconis ICM talk in 1990

Ron Graham ICM talk in 1983



William M. Kantor ICM talk in 1998

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Card Shuffle Groups

The classification of $\langle O, I \rangle$ in [1]

Size of each pile <i>n</i>	$\langle O, I \rangle$
$n = 2^{f}$ for some positive integer f	$C_2 \wr C_{f+1}$
$n \equiv 0 \pmod{4}$, $n > 12$ and n is not a power of 2	$C_2^{n-1} \rtimes A_n$
$n \equiv 1 \pmod{4}$	$C_2^n \rtimes A_n$
$n\equiv 2 \pmod{4}$ and $n>6$	$C_2 \wr \operatorname{Sym}(n)$
$n \equiv 3 \pmod{4}$	$C_2^{n-1} \rtimes S_n$
<i>n</i> = 6	$C_2^6 \rtimes \mathrm{PGL}(2,5)$
<i>n</i> = 12	$C_2^{11} \rtimes M_{12}$

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Card Shuffle Groups

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A deck of *kn* cards with $k \ge 2$

• cut into k piles and then perfectly interleave them (k! ways).

$$0 \qquad 0 \qquad n \qquad \cdots \qquad (k-1)n$$

$$1 \qquad 1 \qquad 1+n \qquad \cdots \qquad 1+(k-1)n$$

$$\vdots \qquad \vdots \qquad \vdots \qquad \vdots \qquad \vdots$$

$$kn-1 \qquad n-1 \qquad 2n-1 \qquad \cdots \qquad kn-1$$

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- ρ_{τ} : the permutation of the *kn* cards induced by the permutation τ of the *k* piles.

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- the *i*th row of the *j*th column is the i + jnth position,

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- the *i*th row of the *j*th column is the *i* + *jn*th position,
- recall σ and $\rho_{\tau} \implies (i + jn)^{\sigma} = ik + j$ and $(i + jn)^{\rho_{\tau}} = i + j^{\tau}n$.

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- The shuffle group on kn cards, denoted by $G_{k,kn}$, is generated by all possible shuffles $\rho_{\tau}\sigma$ for $\tau \in \text{Sym}(\{0, \ldots, k-1\})$.

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- The shuffle group on kn cards, denoted by G_{k,kn}, is generated by all possible shuffles ρ_τσ for τ ∈ Sym({0,..., k − 1}).
 (G_{k,kn} = ⟨ρ_τσ | τ ∈ Sym(k)⟩ = ⟨ρ_τ, σ | τ ∈ Sym(k)⟩.)

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 - $G_{3,3n} \ge \operatorname{Alt}(3n)$ if *n* is not a power of 3;

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- In [2] they also proved:
 - ► $G_{k,kn} \leq \operatorname{Alt}(kn)$ if and only if either $n \equiv 0 \pmod{4}$, or $n \equiv 2 \pmod{4}$ and $k \equiv 0$ or 1 (mod 4).
 - $G_{k,k^m} = \operatorname{Sym}(k) \wr C_m$.

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• Cohen, Harmse, Morrison and Wright^[3] confirmed the latter part of MM's conjecture when k = 4.

 $(G_{4,2^m} = AGL(m, 2) \text{ for some odd integer } m \geq 3)$

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• In [3] they also posed:

Shuffle Group Conjecture (2005)

For $k \ge 3$, if *n* is not a power of *k* and $(k, n) \ne (4, 2^{f})$ for any positive integer *f*, then $G_{k,kn} \ge A_{kn}$.

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 - ▶ k > n;
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 - ▶ k is a power of 2.
- In [4] they also opened up the study of "generalized shuffle groups".

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Our contribution

Theorem (Xia-Zhang-Z. 2023⁺)

The Shuffle Group Conjecture holds when $k \ge 4$ or k does not divide n.
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- Reduction Lemma: If $G_{k,kn}$ is 2-transitive, then either k = 4 and n is a power of 2, or $G_{k,kn}$ contains A_{kn} .
- 2-transitivity Lemma: If either $k \ge 4$ and n is not a power of k, or k = 3 and n is not divisible by 3, then $G_{k,kn}$ is 2-transitive.

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The Shuffle Group Conjecture holds when $k \ge 4$ or k does not divide n.

We established two key lemmas to prove the theorem.

- Reduction Lemma: If $G_{k,kn}$ is 2-transitive, then either k = 4 and n is a power of 2, or $G_{k,kn}$ contains A_{kn} .
- 2-transitivity Lemma: If either $k \ge 4$ and n is not a power of k, or k = 3 and n is not divisible by 3, then $G_{k,kn}$ is 2-transitive.
- A permutation group G on a set Ω is said to be 2-transitive if the induced action of G on Ω × Ω \ {(α, α) | α ∈ Ω} is transitive.

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- The fixed point ratio of a permutation g on a finite set Ω, denoted by fpr(g), is defined as fpr(g) = |Fix(g)|/|Ω|, where Fix(g) = {α ∈ Ω | α^g = α}.
- Observation: $\operatorname{fpr}(\rho_{\tau}) = \operatorname{fpr}(\tau) \Longrightarrow \operatorname{fpr}(\rho_{\tau}) = (k-2)/k$ when τ is a transposition.

Write $G = G_{k,kn}$. Suppose G is an affine 2-transitive permutation group, i.e., $G \leq AGL(V) = AGL(d, p)$ for some d-dimension vector space V over a prime field \mathbb{F}_p .

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- $kn = |V| = p^d \implies (k, n)$ is either $(3, 3^{d-1})$ or $(4, 2^{d-2})$.
- Note that $G_{3,3^d}$ and $G_{4,2^d}$ have been determined.

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- Shuffle Group Conjecture $\xrightarrow{Reduction Lemma}$ 2-transitivity of $G_{k,kn}$.
- Thus the remaining unresolved case of the conjecture is that k = 3 divides *n*.

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- we complete this proof through the following cases:
 - $k \nmid n$ with $k \ge 4$; (only elaborate on this case here)
 - $k \nmid n$ with k = 3;
 - $k \mid n$ with $k \geq 4$.

0	п	•	•	•	(k-1)n
1	1 + n	•	•	•	1+(k-1)n
•	•				•
•	•				•
•	•				•
n-1	2 <i>n</i> – 1		•		kn-1

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Thank you for listening!