# Card Shuffle Groups 

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This is joint work with Binzhou Xia and Zhishuo Zhang.
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## Perfectly shuffle $2 n$ cards

- Cut the deck in half:

- Perfectly interleave them:



## Questions

Perform out-shuffles on a deck of 52 cards repeatedly.

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Perform out-shuffles on a deck of 52 cards repeatedly.

Question: Can it return to the original order?
Answer: Yes. For example, after 52! times.
Question: What is the minimum number of times needed to return to the original order?

Answer: 8 times.

## Why 8 times

- Position $x: 0 \begin{array}{llllllllll} & 1 & 2 & 3 & 4 & \cdots & 25 & 26 & \cdots & 50 \\ 51\end{array}$


Out-shuffle $O$

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 after $O: 0 \begin{array}{llllllllll} & 26 & 1 & 27 & 2 & \cdots & 38 & 13 & \cdots & 25\end{array} 51$


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- $(i+26 j)^{O}=2 i+j$ for $i \in\{0, \ldots, 25\}$ and $j \in\{0,1\}$;


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- $x^{O}=(2 \ell \bmod 2 n-1)$ for $x \in\{1, \ldots, 2 n-2\}$.


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Question: Is it possible to send a given card to any chosen position by performing a sequence of the two shuffles?

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Question: How many/Which orderings can be obtained by performing a sequence of the two shuffles?

To answer these questions, we first determine the parity of $O$ and $I$.

## Observation

Observation: for $n=4$,

- Position $x: 0 \begin{array}{llllllll}1 & 2 & 3 & 4 & 5 & 6 & 7\end{array}$
after $O: 0$
$x^{O}: 0$ 2


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after $O: 0 \begin{array}{llllllll}0 & 4 & 1 & 5 & 2 & 6 & 3 & 7\end{array}$
$x^{0}: 0 \begin{array}{llllllll} & 2 & 4 & 6 & 1 & 3 & 5 & 7\end{array}$
- the inversion number of $(0,2,4,6,1,3,5,7)$ is $1+2+3$.


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| ---: | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 7 |  |  |  |  |  |  |  |
| $x^{O}:$ | 0 | 2 | 4 | 6 | 1 | 3 | 5 |
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For general $n$, the order of the $2 n$ cards after $O$ is

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and thus its inversion number is $1+\cdots+n-1=n(n-1) / 2$.

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\left(x^{\prime}=x^{(0, n)(1, n+1) \cdots(n-1,2 n-1) O} \text { for all } x \in\{0,1, \ldots, 2 n-1\}\right)
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- The permutation of the $2 n$ cards induced by permutating the two piles has the same parity as $n$.
- If $n$ and $O$ have the same parity, then $I$ is even; otherwise $I$ is odd $\Longrightarrow$ If $n \equiv 0$ or $3(\bmod 4)$, then $I$ is even; otherwise $I$ is odd.


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- The permutation of the $2 n$ cards induced by permutating the two piles has the same parity as $n$.
- If $n$ and $O$ have the same parity, then $I$ is even; otherwise $I$ is odd $\Longrightarrow$ If $n \equiv 0$ or $3(\bmod 4)$, then $l$ is even; otherwise $l$ is odd.
- Thus $\langle O, I\rangle \leq \operatorname{Alt}(2 n) \Longleftrightarrow n \equiv 0(\bmod 4)$.


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Answer: Both no.

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- out-shuffle and in-shuffle preserve the partition $\{0,7\},\{1,6\},\{2,5\},\{3,4\}$.

For a general $n$, out-shuffle and in-shuffle preserve the partition $\{0,2 n-1\},\{1,2 n-2\}, \ldots,\{n-1, n\}$.

## $\langle O, I\rangle$ is neither $\operatorname{Alt}(2 n)$ nor $\operatorname{Sym}(2 n)$

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- $O$ and $I$ preserve $\{0,2 n-1\},\{1,2 n-2\}, \ldots,\{n-1, n\}$
$\Longrightarrow\langle O, I\rangle$ is imprimitive.
- $\operatorname{Alt}(2 n)$ is primitive $\Longrightarrow \operatorname{Alt}(2 n) \notin\langle O, I\rangle \Longrightarrow O$ and $I$ can't generate $\operatorname{Alt}(2 n)$ or $\operatorname{Sym}(2 n)$.


## Diaconis-Graham-Kantor

Question: How many/Which orderings can be obtained by performing a sequence of the two shuffles?
[1] P. Diaconis, R. L. Graham and W. M. Kantor, The mathematics of perfect shuffles., Adv. Appl. Math., 4 (1983), 175-196.

## Diaconis-Graham-Kantor

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Answered by Diaconis, Graham and Kantor in 1983 ${ }^{[1]}$.


Persi Diaconis
ICM talk in 1990


Ron Graham
ICM talk in 1983


William M. Kantor
ICM talk in 1998
[1] P. Diaconis, R. L. Graham and W. M. Kantor, The mathematics of perfect shuffles., Adv. Appl. Math., 4 (1983), 175-196.

## The classification of $\langle O, I\rangle$ in [1]

| Size of each pile $n$ | $\langle O, I\rangle$ |
| :--- | :--- |
| $n=2^{f}$ for some positive integer $f$ | $C_{2} \prec C_{f+1}$ |
| $n \equiv 0(\bmod 4), n>12$ and $n$ is not a power of 2 | $C_{2}^{n-1} \rtimes A_{n}$ |
| $n \equiv 1(\bmod 4)$ | $C_{2}^{n} \rtimes A_{n}$ |
| $n \equiv 2(\bmod 4)$ and $n>6$ | $C_{2} \imath \operatorname{Sym}(n)$ |
| $n \equiv 3(\bmod 4)$ | $C_{2}^{n-1} \rtimes S_{n}$ |
| $n=6$ | $C_{2}^{6} \rtimes \operatorname{PGL}(2,5)$ |
| $n=12$ | $C_{2}^{11} \rtimes M_{12}$ |

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## A deck of $k n$ cards with $k \geq 2$

- cut into $k$ piles and then perfectly interleave them ( $k$ ! ways).

| 0 |  | 0 | $n$ | $\cdots$ | $(k-1) n$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 |  | 1 | $1+n$ | $\cdots$ | $1+(k-1) n$ |
| $\vdots$ | $\longrightarrow$ | $\vdots$ | $\vdots$ |  | $\vdots$ |
| $k n-1$ |  | $\vdots$ | $\vdots$ |  | $\vdots$ |
| $\vdots$ |  | $n-1$ | $2 n-1$ | $\cdots$ | $k n-1$ |

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| - | - | - |  | - |
| - | - | - |  | - |
| $k n-1$ | $n-1$ | $2 n-1$ | -•• | kn-1 |

- Standard shuffle $\sigma$ : picking up the top card from each of the piles $0, \ldots, k-1$ in order and repeating until all cards have been picked up.


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- $\rho_{\tau}$ : the permutation of the $k n$ cards induced by the permutation $\tau$ of the $k$ piles.


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- For all $i \in[n]$ and $j \in[k]$,
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- the $i$ th row of the $j$ th column is the $i+j n$th position,
- recall $\sigma$ and $\rho_{\tau} \Longrightarrow(i+j n)^{\sigma}=i k+j$ and $(i+j n)^{\rho_{\tau}}=i+j^{\tau} n$.


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- The shuffle group on $k n$ cards, denoted by $G_{k, k n}$, is generated by all possible shuffles $\rho_{\tau} \sigma$ for $\tau \in \operatorname{Sym}(\{0, \ldots, k-1\})$.


## Shuffle groups

- For a positive integer $m$, denote $[m]=\{0,1, \ldots, m-1\}$.
- $k$ piles:

| 0 | $n$ | $\cdots$ | $(k-1) n$ |
| :---: | :---: | :---: | :---: |
| 1 | $1+n$ | $\cdots$ | $1+(k-1) n$ |
| $\vdots$ | $\vdots$ |  | $\vdots$ |
| $n-1$ | $2 n-1$ | $\cdots$ | $k n-1$ |

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$$
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- In [2] they also proved:
- $G_{k, k n} \leq \operatorname{Alt}(k n)$ if and only if either $n \equiv 0(\bmod 4)$, or $n \equiv 2$ $(\bmod 4)$ and $k \equiv 0$ or $1(\bmod 4)$.
- $G_{k, k^{m}}=\operatorname{Sym}(k) \imath C_{m}$.
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- In [3] they also posed:


## Shuffle Group Conjecture (2005)

For $k \geq 3$, if $n$ is not a power of $k$ and $(k, n) \neq\left(4,2^{f}\right)$ for any positive integer $f$, then $G_{k, k n} \geq A_{k n}$.

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- In [4] they also opened up the study of "generalized shuffle groups".

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## Our contribution

Theorem (Xia-Zhang-Z. 2023+ ${ }^{+}$)
The Shuffle Group Conjecture holds when $k \geq 4$ or $k$ does not divide $n$.

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We established two key lemmas to prove the theorem.

- Reduction Lemma: If $G_{k, k n}$ is 2-transitive, then either $k=4$ and $n$ is a power of 2 , or $G_{k, k n}$ contains $A_{k n}$.
- 2-transitivity Lemma: If either $k \geq 4$ and $n$ is not a power of $k$, or $k=3$ and $n$ is not divisible by 3 , then $G_{k, k n}$ is 2-transitive.


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- A permutation group $G$ on a set $\Omega$ is said to be 2-transitive if the induced action of $G$ on $\Omega \times \Omega \backslash\{(\alpha, \alpha) \mid \alpha \in \Omega\}$ is transitive.


## Reducing to 2-transitivity

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- The fixed point ratio of a permutation $g$ on a finite set $\Omega$, denoted by $\operatorname{fpr}(g)$, is defined as $\operatorname{fpr}(g)=|\operatorname{Fix}(g)| /|\Omega|$, where $\operatorname{Fix}(g)=\left\{\alpha \in \Omega \mid \alpha^{g}=\alpha\right\}$.


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- Observation: $\operatorname{fpr}\left(\rho_{\tau}\right)=\operatorname{fpr}(\tau) \Longrightarrow \operatorname{fpr}\left(\rho_{\tau}\right)=(k-2) / k$ when $\tau$ is a transposition.


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Write $G=G_{k, k n}$. Suppose $G$ is an affine 2-transitive permutation group, i.e., $G \leq \operatorname{AGL}(V)=\operatorname{AGL}(d, p)$ for some $d$-dimension vector space $V$ over a prime field $\mathbb{F}_{p}$.

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- Note that $G_{3,3^{d}}$ and $G_{4,2^{d}}$ have been determined.


## 2-transitivity of $G_{k, k n}$

2-transitivity Lemma (Xia-Zhang-Z. 2023 ${ }^{+}$)
If either $k \geq 4$ and $n$ is not a power of $k$, or $k=3$ and $n$ is not divisible by 3 , then $G_{k, k n}$ is 2 -transitive.

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- Shuffle Group Conjecture $\xrightarrow{\text { Reduction Lemma }}$ 2-transitivity of $G_{k, k n}$.
- Thus the remaining unresolved case of the conjecture is that $k=3$ divides $n$.


## Sketch of Proof to 2-transitivity Lemma

- Let $G=G_{k, k n}$ and $[m]=\{0,1, \ldots, m-1\}$.


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- we complete this proof through the following cases:
- $k \nmid n$ with $k \geq 4$; (only elaborate on this case here)
- $k \nmid n$ with $k=3$;
- $k \mid n$ with $k \geq 4$.


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- $x^{\sigma}=k x>x \Longrightarrow x \in n^{G_{0}}$.


## Thank you for listening!


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